

# A Secure Mobile RFID System Design

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**Abstract:** Mobile RFID is a new application that provides a valuable service to users with mobile communication. Current RFID authentication schemes are not viable to use in mobile RFID environment. Since unlike in regular RFID environments, whereas the communication channel between the server and the reader is assumed to be secure, the communication channel between the server and reader in the mobile RFID system is not assumed to be secure. Thus, it need to study a new communication system that provides a secure privacy of mobile RFID devices. Therefore, this paper proposes a mobile RFID authentication system to provide a privacy mechanism updating confidential information in order to protect many security and privacy threats for mobile RFID devices.

**Keywords:** Mobile RFID, Privacy, Authentication Protocol, Process Domain, Node Domain, Network Domain.

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## Introduction

Radio Frequency Identification (RFID) has been widely used in many applications and is an automatic identification technology that allows remote interrogation of the ID data on RFID tags using radio frequency in accordance with wireless communication between tagged objects and RFID readers[1- 9].

Privacy preserving authentication in RFID systems is an important problem. In secure RFID systems, a reader will accept the tag's information only after it authenticates a tag. The tag may be attached to a pharmaceutical product that carries patient information, a passport that carries a personal identification, or a commercial product that carries information about manufactured date, expiring date, ingredients, etc.

Ordinary RFID is believed that the communication channel between the reader and the database is safe. However, in the mobile RFID system, the communication channel between the reader and the database is using wireless channel, and it is not assumed to be safe. The insecure channel is vulnerable to various threats such as eavesdropping, business espionage, and tag masquerading. For this reason, the mobile RFID system has security and privacy problems are difficult and more challenging than those of order RFID system.

Recently, some authentication protocol [10-13] were proposed for secure mobile RFID systems. The proposed protocols either need additional devices, or require complex encryption and decryption calculations. They have a common inherent problem that parties in their protocol must be to share one or more secret key.

Almost all of the existing mobile RFID authentication protocol is followed traditional challenge-response mechanism. The mobile reader is sending a random challenge to the tag, then the tag is replying with its authentication message which be generated by its secret value and other random numbers. An adversary can transmit intended or meaningless requests to acquire the location key and identification information from the tag.

The adversary can anticipate the response message of the tag and can use it to perform location tracking. The tag information can be obtained by sending some intended requests.

Several schemes [14-16] has been proposed, which use only ultra-light weight, such as the random number generator, pseudo random number generator, cyclic redundancy check and exclusive-or operation, to resist the latent attacks. However, these schemes still suffer from security weakness because the schemes apply only ultra-light weight operations.

The rest of this paper is organized as follows: related work is presented in section 2, authentication system for mobile RFID device is presented in section 3, analysis is detailed in section 4, and section 5 concludes the paper.

This work was supported by the National Research Foundation of Korea(NRF) and the Center for Women In Science, Engineering and Technology(WISET) Grant funded by the Korean Government(Program for Returners into R&D by the Ministry of Science, ICT & Future Planning(MSIP)). This work was partly supported by the ICT R&D program of MSIP/IITP, Republic of Korea. [13-912-06-003, Development of Mobile S/W Security Testing Tools for Detecting New Vulnerabilities of Android]

### Related Works

The mobile RFID-based applications lately become a new issue for research and development due to advantages of RFID technology and mobile smart device. However, a mobile RFID system is encountering all kinds of security and privacy threats because the communication between a reader and a database is insecure.

[17] scheme was also proposed to get rid of some security issues which are nothing but an extra device added approach. [18] scheme introduced a prototype named “watchdog tag” and [19] scheme proposed two protocols to provide low-cost RFID systems that require good operational and security properties. The first protocol of [19], in spite of very restricted functionality, resolves not only security properties, but also operational properties, the second protocol of [19] resolves the replay attack in the cost of a random number generator in tags. Moreover, these schemes have significantly reduced the amount of tag transmissions which is the most energy consuming task. [20] scheme proposed a low-cost RFID mutual authentication protocol based on the method of Hash-based Message Authentication Code (HMAC) under the assumption that hash function is secure, the property that the new protocol can achieve mutual authentication between reader and tag. [21] scheme proposed a one-way hash based low-cost authentication protocol with forward security and analyzed its efficiency, but the computation load was not taken into consideration. [22] scheme improved two lightweight RFID authentication protocols to pursue stronger anonymity property and security feature. [23] analyzed and improved lightweight RFID protocol using substring. [24] scheme proposed an effective and secured certificate mechanism using mobile devices as RFID readers together with the credit cards containing RFID tags. The result shows it can improve the existing RFID security issue under the premise of safety, efficiency and compatibility. [25] scheme proposed ultra-lightweight RFID authentication. However, the schemes did not have a mechanism to verify updated secret key between the tag and the server.

Meanwhile, a mobile phone embedded RFID reader modules will be situated anywhere and operated using many RFID tags in various RFID application systems. It is difficult to secure the privacy of a mobile RFID-enabled device without an authentication system with a novel communication protocol. Therefore, It is necessary for a new authentication system to protect against privacy threats to mobile RFID readers.

### Authentication System for Mobile RFID Device

Currently, RFID authentication research cannot be used in mobile RFID environment and the authentication results in computing load to low cost tag. Therefore, to provide sufficient protection to the information privacy, this paper proposes three-layer mobile RFID authentication system for privacy protection.

#### A. Layer-based Mobile RFID Module

Mobile RFID system has network, process and node domain as shown in Fig.1. The mechanism constitutes functional layers for an efficient identification and each layer supports mobile RFID authentication module.

Table1 shows Three-layer of network: Network domain has group tag  $G_i$  and group tag size  $G_{size}$ . Node domain contains reader node  $N_R$  and tag node  $N_T$ . Process domain includes the authentication of reader, tag and server etc.

Table 1. Notation

3-Layer Domain Name	3-Layer Contents
Network Domain	$G_i, G_{size}$
Node Domain	$N_R, N_T$
Process Domain	Authentications in Reader, Tag, Server etc.

Network domain constitutes a group tag  $G_i$  and group tag size  $G_{size}$ .  $G_i$  denotes a group classification compliant with a purpose of tag use and  $G_{size}$  denotes the number of tag in a group. If some  $G_i$  are same purpose,  $G_{size}$  starts query by size. Network domain provides quick and correct response from the tag query by connecting the groups from a purpose of tag use.

Node domain constitutes  $N_R$  and  $N_T$  to verify updated secret keys.  $N_R$  acts to manage shared secret keys between a reader and a server and  $N_T$  acts to manage shared secret keys between a tag and a server. After setting up  $ID_{deadline}$ , node domain informs  $N_R$  about  $K_i$  secret key shared between a reader and a server. Similarly, it informs  $N_T$  about  $K_i$  secret key shared between a tag and a server. In addition, it verify forward secrecy for the updated secret keys of  $N_R$  and  $N_T$ .

Process domain as a core mobile RFID authentication system proceeds authentication protocol in a reader, a tag and a server.

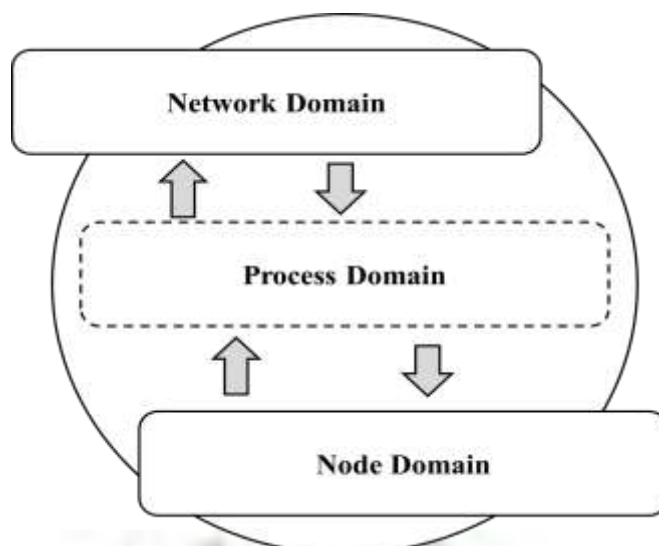


Fig. 1. Three-layer based a Mobile RFID Mechanism

In Fig.1, after process domain interacts with network domain, it can completely verify a mobile RFID system by means of node domain.

### B. Mobile RFID System Authentication Phase

The proposed authentication module in Fig.2 circulates such as initialization, registration, tag authentication, secret key updating in a server, encryption, reader authentication, secret key updating in a tag and verification of updated secret key. The authentication phase supports lightweight authentication to be implemented on low-cost tags that ensure privacy protection.



Figure 2. Mobile RFID System Authentication Phase

### C. Mobile RFID Authentication Protocol

Table 2 shows the notations for the proposed authentication protocol.

Table 2. Notation

Symbol	Description
$ID_T$	Identification of the Tag
$ID_R$	Identification of the Reader
$ID_{deadline}$	Identification Deadline
$K_R$	Reader Password
$D$	Detailed Information about the Tag in the Database
$K_i$	Secret Key shared between the tag and the Server
$K_{i+1}$	Updated Secret Key used in between the Tag and the Server
$s$	Secret Key shared between the Reader and the Server
$r$	A random number
$G_t, G_{size}$	Group Tag, Group Tag Size
$N_R, N_T$	Reader Node, Tag Node
$G()$	Generator function
$H()$	Hash function
$\oplus$	XOR
$\parallel$	Concatenation
$TSP$	Timestamp

The mobile RFID reader must be registered and authenticated by the server. After the server has authenticated the reader, it sends  $ID_R$  and  $K_R$  to the reader. Fig.3 explains the mobile RFID authentication protocol as follows.

#### Initialization Phase:

1. The reader generates and stores a random number  $r$  through the use of a  $G()$  within the reader
2. The tag confirms its use purpose by sending a query to  $G_t$  of network domain.

#### Registration Phase:

3. The reader sends  $H(ID_R)$  and  $H(ID_T \oplus K_i)$  to the server.
4. After receiving the purpose message of tag use from  $G_t$  of network domain, the tag sends  $H(ID_T \oplus K_i)$  to the reader.

#### Tag Authentication Phase:

5. The server checks whether  $H(ID_T \oplus K_i)$  matches with the stored hash value of the tags. If it matches, the tag is authenticated else failed.

#### Update Secret Key in the Server Phase:

6. The server updates  $K_i$  to  $K_{i+1}$  where  $K_{i+1} = G(K_i)$
7. The server sends  $K_{i+1}$  and  $s$  to  $N_R$  of node domain

#### Encryption Phase:

8. The server computes  $H(ID_T \oplus K_i) \oplus K_{i+1}$ .
9. The tag sends  $E_{KR}[D \parallel H(ID_T \oplus K_i) \oplus K_{i+1}]$  with  $K_R$  to the reader.
10. After decrypting it, the reader takes  $D$  and sends  $H(ID_T \oplus K_i) \oplus K_{i+1} \oplus r \oplus TSP$  to the tag.



#### Reader Authentication Phase:

11. The tag verifies the reader by  $r$  and  $TSP$ , then it computes  $K_{i+1}$  of the updated secret key between the tag and the server using  $H(ID_T \oplus K_i)$  with  $r$ .

#### Update Secret Key in the Tag Phase:

12. The tag updates  $K_i$  to  $K_{i+1}$ .
13. The tag sends  $K_{i+1}$  to  $N_T$  of node domain

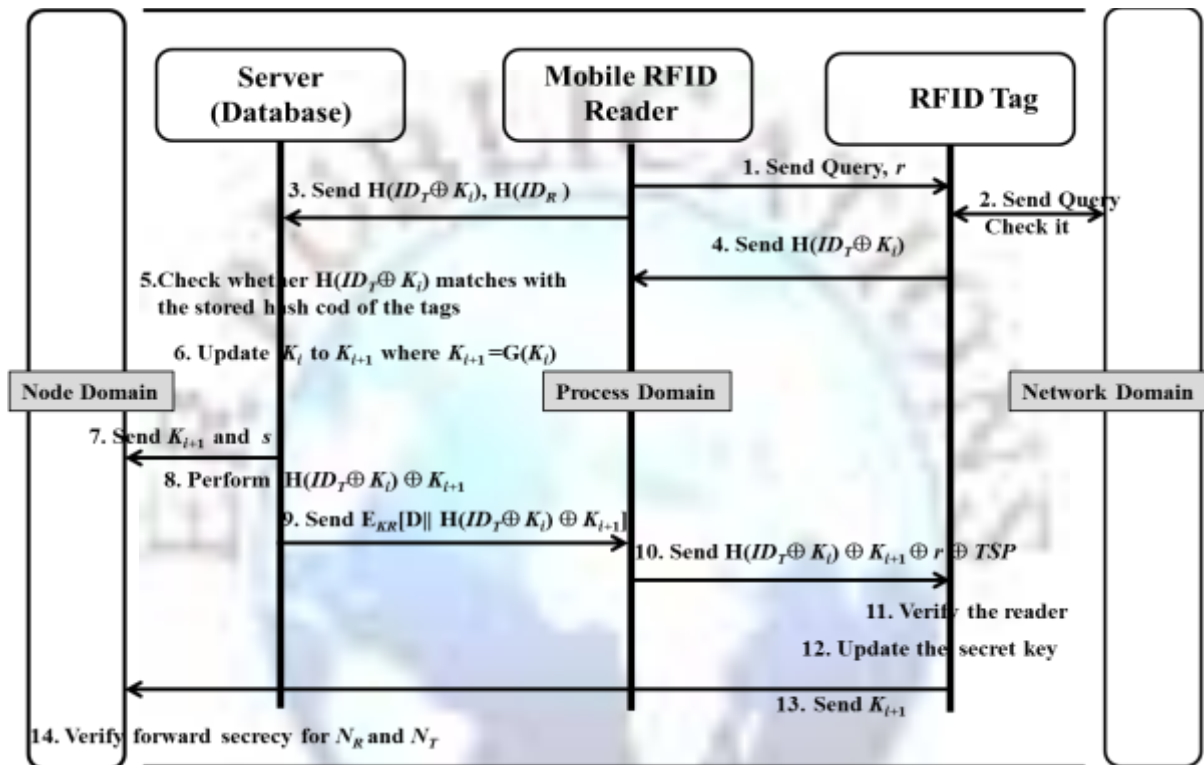


Figure 3. Mobile RFID Authentication Protocol

#### Verification for Update Secret Key Phase:

14. Node domain verifies that the updated secret keys of  $N_R$  and  $N_T$  are carried out within  $ID_{deadline}$ .  
Updating mechanism for confidential information conforms to node domain. The server which stores old and new values of the tag secrecy communicates with node domain and node domain manages the updated secret keys. Using forward secrecy, the node domain verifies the updated secret keys of  $N_R$  and  $N_T$  within  $ID_{deadline}$ .

#### Analysis

In this section, the computation cost and security strength of the proposed method is analyzed. Table 3 compares the proposed scheme with the related schemes in terms of the computation cost during the authentication phase. In Table 3,  $T_{XOR}$  is the execution time and the computation cost for the proposed scheme has lower computation-cost operation than other schemes.

Table 3. Notation

Schemes	Computation Costs	
	Tag	Server
Chien and Chen[14]	$2T_{XOR}$	$1T_{XOR}$
Chen and Deng[15]	$5T_{XOR}$	$3T_{XOR}$
Huang and Jiang[16]	$6T_{XOR}$	$6T_{XOR}$
Proposed Scheme	$1T_{XOR}$	$1T_{XOR}$

Table 4 shows the comparisons of schemes in terms of what attacks they can resist. Huang and Jiang's scheme cannot resist the tracking attack. Chien and Chen's scheme and Chen and Deng's scheme suffer from all the attacks. However, the proposed scheme can resist all the attacks.

Table 4. Security Comparisons

Schemes Attacks	Chien and Chen[14]	Chen and Deng[15]	Huang and Jiang[16]	Proposed Scheme
Resist forged-server	No	No	Yes	Yes
Resist forged-tag	No	No	Yes	Yes
Resist tracking	No	No	No	Yes
Resist replay	No	No	Yes	Yes
Resist forward secrecy	No	No	Yes	Yes
Resist DoS	No	No	Yes	Yes

The security of the proposed scheme against relevant attacks is analyzed as follows:

#### 1) Replay Attack

The attackers can obtain outputs of the tag and transmit the eavesdropped messages to the reader. However, he cannot impersonate the legitimate tag because the outputs are different on every session due to random number  $r$  and  $TSP$

#### 2) Forward Security

The forward security means that even if the adversary obtains a current secret key, he still cannot derive the keys used for past time periods. To ensure this, verification for update secret key scheme which involves key deployment is used. For each valid session, a tag uses the current key  $K_i$  for generation and verification of authentication tags. At the end of each valid session,  $K_i$  is updated by generator function  $G(\cdot)$  and previous  $K_i$  is deleted. After a server sends  $K_{i+1}=G(K_i)$  to  $N_R$ ,  $N_R$  confirms and manages that the current key  $K_i$  updates for forward security within identification deadline  $ID_{deadline}$  at the end session. Similarly,  $N_T$  confirms and manages it. Therefore, even if the adversary obtains the current secret key, he cannot induce any previous keys due to the protection of updated secret key from node domain.

#### 3) Privacy Attack

An attacker in privacy attack wants to know the query and information of a tag. In initialization and registration phase of paragraph 3.3, the query and information of a tag is protected. In addition, an attacker cannot know the query and information of a tag because the tag sends  $E_{KR}[D||H(ID_T \oplus K_i) \oplus K_{i+1}]$  with a reader password  $K_R$  to the reader in encryption phase.

#### 4) Tag Cloning

Tag cloning means that the data on a valid tag is scanned and copied by a malicious RFID reader and the copied data is embedded onto a fake tag. After tag authentication and updating secret key in the server, the encryption and reader authentication by  $r$  and  $TSP$  are carried out. These mean that the tag cannot respond true values unless the tag first authenticates the reader. Moreover, they prevent a malicious reader from scanning and copying the data on a valid tag.

**5) Eavesdropping**

The information of proposed methods has been encoded by hash function which makes the adversary not to get the true value because of the one-way characteristic. The attackers cannot know the detailed content of the information even if they spy on the outputs. In encryption phase, the tag sends its detail in encrypted form with a reader password  $K_R$  to the reader, so the attackers cannot know the real information.

**6) Denial of Service Attack**

The proposed scheme needs synchronization of secret key shared between the server and the tag. In tag authentication phase, after the server checks whether  $H(ID_T \oplus K_i)$  matches with the stored hash value of the tags, it updates  $K_i$  to  $K_{i+1}$ . If the adversary performs this attack on the flow the proposed protocol, he can prevent the tag from taking confirmation. This breaks the synchronization between the tag and the server because the server refreshes the tag secrecy but the tag does not. However, in the proposed protocol, the server makes itself synchronize with the tag at the situation because it stores old and new values of the tag secrecy.

**Acknowledgment**

The authors appreciate Prof. Cheong Youn and Eunbae Kong for their helpful research supporting.

**Conclusion**

The mobile RFID applications are becoming a new issue of their research and development because of the advantages of RFID technology and mobile smart devices. However, the mobile RFID system encounters various problems of security and privacy because the communication channel between the server and the reader in the mobile RFID system is not secure. If the security problems is not solved, the use of mobile RFID applications will be greatly reduced as well as the mobile RFID applications are threatened.

Therefore, the proposed scheme solves the problem of security and privacy from the unsecure communication channel between the server and the reader because of the encryption using hash functions satisfying the characteristics of the mobile RFID system such as low computation load. To discourage fraud of valid mobile reader, the updating mechanism for secret information can provide the privacy of a mobile RFID system using authentication of the tag and the reader. In addition, the proposed system verifies the updated secret key from the server, so its security is stronger than previous mobile RFID system.

As a result, the proposed system demonstrates its superiority to other related schemes in terms of computation cost and security.

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